



**AB-34**

**APPLICATION  
BRIEF**

**Integer Square Root Routine  
for the 8096**

**LIONEL SMITH  
ECO APPLICATIONS ENGINEER**

April 1989



Order Number: 270523-001

Information in this document is provided in connection with Intel products. Intel assumes no liability whatsoever, including infringement of any patent or copyright, for sale and use of Intel products except as provided in Intel's Terms and Conditions of Sale for such products.

Intel retains the right to make changes to these specifications at any time, without notice. Microcomputer Products may have minor variations to this specification known as errata.

\*Other brands and names are the property of their respective owners.

†Since publication of documents referenced in this document, registration of the Pentium, OverDrive and iCOMP trademarks has been issued to Intel Corporation.

Contact your local Intel sales office or your distributor to obtain the latest specifications before placing your product order.

Copies of documents which have an ordering number and are referenced in this document, or other Intel literature, may be obtained from:

Intel Corporation  
P.O. Box 7641  
Mt. Prospect, IL 60056-7641  
or call 1-800-879-4683

**INTEGER SQUARE ROOT  
ROUTINE FOR THE 8096**

**CONTENTS**

PAGE

Theory ..... 1  
Practice ..... 1  
Comments ..... 1





This Application Brief presents an example of calculating the square root of a 32-bit signed integer.

## Theory

Newton showed that the square root can be calculated by repeating the approximation:

$$X_{new} = (R/X_{old} + X_{old})/2 ; X_{old} = X_{new}$$

where: R is the radicand

Xold is the current approximation of the square root

Xnew is the new approximation

until you get an answer you like. A common technique for deciding whether or not you like the answer is to loop on the approximation until Xnew stops changing. If you are dealing with real (floating point) numbers this technique can sometimes get you in trouble because it's possible to hang up in the loop with Xnew alternating between two values. This is not the case with integers. As an example of how it all works, consider taking the square root of 37 with an initial guess (Xold) of 1:

$$X_{new} = (37/1 + 1)/2 = 19; X_{old} = 19$$

$$X_{new} = (37/19 + 19)/2 = 10; X_{old} = 10$$

$$X_{new} = (37/10 + 10)/2 = 6; X_{old} = 6$$

$$X_{new} = (37/6 + 6)/2 = 6; X_{old} = 6 - \text{done!}$$

Note that in integer arithmetic the remainder of a division is ignored and the square root of a number is floored (i.e. the square root is the largest integer which, when multiplied by itself, is less than or equal to the radicand).

## Practice

The only significant problem in implementing the square root calculation using this algorithm is that the division of R by Xold could easily be a 32 by 32 divide if R is a 32 bit integer. This is ok if you happen to have a 32 by 32 divide instruction, but most 16-bit machines (including the 8096) only provide a 32 by 16 divide. However, a little bit of creative laziness will allow us to squeeze by using the 32 by 16 bit divide on the 8096.

The largest positive integer you can represent with a 32-bit two's complement number is 07fff\$ffffh, or 2,147,483,647. The square root of this number is 0b504h, or 46,340. The largest square root that we can generate from a 32-bit radicand can be represented in 16-bits. If we are careful in picking our initial Xold we can do all of the divisions with the 32 by 16 divide instruction we have available. Picking the largest possible 16-bit number (0ffffh) will always work although it may slow the calculation down by requiring too many iterations to arrive at the correct result. The algorithm below takes a slightly more intelligent approach. It uses the normalize instruction to figure out how many leading zeros the 32-bit radicand has and picks an initial Xold based on this information. If there are 16 or more leading zeros then the radicand is less than 16 bits so an initial Xold of 0fffh is chosen. If the radicand is more than 16 bits then the initial Xold is calculated by shifting the value 0ffffh by half as many places as there were leading zeros in the radicand. To give credit where credit is due, I first saw this 'trick' in the January 1986 issue of Dr. Dobbs's Journal in a letter from Michael Barr of McGill University.

The routine was timed in a 12.0 Mhz 8096 as it calculated the square roots of all positive 32-bit numbers, the following numbers include the CALL and return sequence and were measured using Timer 1 of the 8096.

Minimum Execution Time: 24 microseconds

Maximum Execution Time: 236 microseconds

Average Execution Time: 102 microseconds

## Comments

The program module which follows is part of a collection of routines which perform integer and real arithmetic on a software implemented tagged stack. The top element of the stack is call TOS and is in fixed locations in the register file. Since the square root operation only involves TOS, further details of the stack structure are not shown.

```

MCS-96 MACRO ASSEMBLER  SQRT                                05/12/86 10:44:30 PAGE 1
DOS MCS-96 MACRO ASSEMBLER, V1.1
SOURCE FILE: ROOT2.A96
OBJECT FILE: ROOT2.OBJ
CONTROLS SPECIFIED IN INVOCATION COMMAND: NOSB
ERR LOC OBJECT      LINE      SOURCE STATEMENT
      1 ;
      2 sqrt module
      3 ;
      4 ; 32 bit integer square root for the 8096
      5 ;
      6 public qstk_isqrt          ; TOP ← SQUARE_ROOT(TOP)
      7 extrn interr:entry        ; Integer error routine
      8 ;
      9 ; id stags for stack integer routines
0019      10 isqrt_id      equ    19h
      11 ;
      12 ; error codes
      13 ;
0000      14 overflow      equ    00h
0001      15 paramerr     equ    01h
0002      16 invalid_input equ    02h
      17
001C      18      oseg at 1ch
      19 ; =====
001C      20 ax:      dsw 1
001C      21 al equ ax:byte
001D      22 ah equ (ax+1):byte
001E      23 dx:      dsw 1
0020      24 cx:      dsw 1
0022      25 bx:      dsw 1
0018      26 sp      equ 18h:word
      27
      28
0030      29      oseg at 30h
      30 ; =====
0030      31 qstk_reg:
0030      32      dsl 1          ; make sure of alignment
0030      33 next      equ qstk_reg:word ; pointer to next element in the arg stack.
0032      34 tos_tag equ (qstk_reg+2):word
0034      35 tos_value:
0034      36      dsl 1          ; 32 bit integer
      37 ;
0000      38      cseg
      39 ; =====
      40 bl macro param
      41      bnc param
      42      endm
      43
      44 bhe macro param
      45      bc param
      46      endm
      47 $eject

```

```

MCS-96 MACRO ASSEMBLER SQRT                                05/12/86 10:44:30 PAGE 2
ERR LOC OBJECT LINE SOURCE STATEMENT
0000 48 cseg
49 ; ====
50 ;
0000 51 qstk_isqrt:
52 ; Takes the square root of the long integer in TOS
53 ; TOS→ Top of argument stack
54 ; iTOS - iSQRT(TOS)
55 ;
0020 56 Xold set cx
0000 A0341C 57 ld ax,tos_value
0003 A0361E 58 ld dx,(tos_value+2)
0006 371F07 59 bbc (dx+1),7,qsi05 ; if (TOS < 0)
0009 C90119 60 push #(isqrt_id*256+paramerr)
000C EF0000 E 61 call interr ; Call interr.
000F F0 62 ret ; Exit
0010 63 qsi05:
0010 0F221C 64 normal ax, bx
0013 DF3E 65 be qstk_isqrt0
0015 991022 66 cmpb bx,#16 ; if (TOS < 2**16)
0018 DA06 67 ble qsil0
001A A1FF0020 68 ld Xold, #0ffh ; Use 0ffh as first estimate.
001E 200A 69 br qstk_isqrtloop
0020 70 qsil0:
0020 180122 71 shrb bx,#1 ; else
0023 A1FFFF20 72 ld Xold, #0ffffh ; Base the first estimate on the
0027 082220 73 shr Xold, bx ; number of leading zeroes in TOS.
002A 74 qstk_isqrtloop;
002A A0341C 75 ld ax,tos_value ; do
002D A0361E 76 ld dx,(tos_value+2) ; if (The divide will overflow)
0030 88201E 77 cmp dx,Xold ; The loop is done.
78 bhe qstk_isqrt_done
0035 8C201C 80 divu ax,Xold ; if ( (ax=TOS/Xold) >= Xold)
0038 88201C 81 cmp ax,Xold ; The loop is done.
82 bhe qstk_isqrt_done
003D 0122 84 clr bx ; Xold=(ax+Xold)/2
003F 641C20 85 add Xold,ax
0042 A40022 86 addc bx,0
0045 0C0120 87 shr1 Xold,#1
0048 27E0 88 br qstk_isqrtloop ; while (The loop is not done)
004A 89 qstk_isqrt_done:
004A A02034 90 ld tos_value,Xold ; TOS=00:Xold
004D A00036 91 ld (tos_value+2),0
0050 92 qstk_isqrt0:
0050 F0 93 ret ; Exit
0051 94 end

```

ASSEMBLY COMPLETED. NO ERROR(S) FOUND.